

# Advances in Compilers

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# Outlines:

- 1 Compiler basics: revisiting
- 2 New vs old compilers
- 3 Parallel language features
- 4 Parallel compilers
- 5 Modern Compilers
- 6 CETUS
- 7 Compiling
- 8 Lexical analyser generator
- 9 Resources for lex

Compilers are language processors (input HLL, output LLL).

- 1 Tokenizing (lexical analysis)
- 2 Parsing (syntax analysis)
- 3 Static semantic checking
- 4 Intermediate code generation
- 5 optimization
- 6 final code generation

# Compiler v/s Interpretation

- Compiled programs run faster, if they are compiled into the form that is directly executable on the underlying hardware
- Static compilation can devote arbitrary amount of time for program analysis and optimization
- Hence JIT (just in time) compiler
- Interpreted programs are typically smaller
- Interpreted programs tends to be more portable

Traditionally there are two approaches to translation:

- 1 Compilation: translates one language to another (say, C to assembly or C to machine). It can be further improved after compilation.
- 2 Interpretation: No more improvement possible, and does immediate execution.

# What practice one should do along with principles?

- Focus on implementing core parts of a compiler, with building the infrastructure
- Small language as an implementing target
- Build from scratch a working interpreter for a small functional language
- The practical should help in understanding the language
- Learning compilers provide strong theoretical foundations.

- Early compilers: Written in low level languages like C or assembly
- Modern compilers are written in C/C++, C#, F#, Java.
- Single individuals usually crafted compilers, but modern compilers are typically large
- In new: there are RISC, CISC machines of past, vector processors, multicore, etc.

- Number of process are running instructions simultaneously
- Instruction cycle?
- Say, fetch time ( $t_f$ ) = decode time ( $t_d$ ) = run-time ( $t_r$ ) =  $t$   $\mu$ -secs, then parallel processing possible.
- Then, a program of 100 machine instructions will take  $300t$   $\mu$  sec. on that machine.
- If instruction  $I_i$  is in execution,  $I_{i+1}$  is in decode cycle, and  $I_{i+2}$  is in fetch cycle. Then, it will take only  $100t$   $\mu$  sec.!!



Consider the following code:

1.  $b = a + c$ ;
2.  $d = b + e$ ;
3.  $f = g + d$ ;

Running 1. and 2. on two cpus at the same time?(Y/N)

Running 1. and 3. on two cpus at the same time?(Y/N)

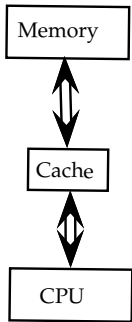
Running bubble sort on two cpus at the same time?(Y/N)

Running quick sort on two cpus at the same time?(Y/N)

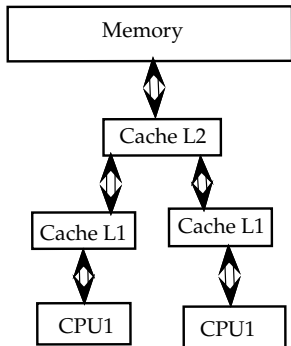
- Compiler algorithms for *parsing*, *type checking*, *data-flow analysis*, *loop transformations* are based on - data dependent analysis in graphs, register allocation based on graph coloring, are new techniques.
- Coding optimization: avoids redundant computations, register allocation, enhancing locality, instruction level parallelism.
- Optimized compilers produce code having some times peak performance of target machine.
- identifying some of the frequent bugs (e.g., improper memory allocation, locations, race conditions, buffer over-run). Thus providing better security.

# Compiler challenges in multicore processors

Single core:



Multicore:



In the era of multicore processors:

- From now on, clock frequencies will rise slowly, if at all,
- but the number of cores on processor is likely to double every couple of years
- by 2020 microprocessors are likely to have number of cores hundreds of even thousands, with heterogeneous functionalities
- Exploiting large -scale parallel will be essential for improving an applications performance
- This is to do without undue programmer effort

The full potential can be achieved only if parallel compilers exist.  
Agenda for the Compiler Community:

- Open compiler architectures,
- Open benchmarks for performance evaluation,
- Make parallel program main-stream
- Write compilers capable of self improvement
- Develop algorithms for optimization of parallel code
- Develop software as reliable as airplane
- Enable system software that is secure at all levels
- Expand compiler courses with new problem domains (such as security)

- Make parallel programming as mainstream
- As on today the parallel programming exists only in databases, and server-side applications
- How to parallel this?:  $\sum_{i=0}^{999} a_i$
- It is normally done as: fetch instruction, fetch data, decode, execute (cycle repeated 1000 times)
- For parallelism fetch opcode once only, decode once only, fetch data 1000 times, run parallel with fetch !!

- Enable system software that is secure at all levels: sophisticated program analysis, prevention of software vulnerabilities, like buffer over-flow and dangling pointers arising from coding defects.
- Automatic verification of complete software stack: “Instead of debugging a program, prove that it meets its specifications, and this proof should be checked by a computer program” - John McCarthy.

# CETUS: A source-to-source compiler infrastructure for multicores

- After the name of a constellation
- Automatic parallelization
- Input C program, output for 64 bit machines
- CETUS symbolic expression tools:

$$1 + 2 * a + 4 - a \Rightarrow 5 + a \text{ (folding)}$$

$$a * (b + c) \Rightarrow a * b + a * c \text{ (distribution)}$$

$$(a * 2) / (8 * c) \Rightarrow a / (4 * c) \text{ (division)}$$

$$(1 - a) < (b + 2) \Rightarrow (1 + a + b) > 0 \text{ (normalization)}$$

$$a \ \&\& \ 0 \ \&\& \ b \Rightarrow 0 \text{ (short-circuit evaluation)}$$



# CETUS: A source-to-source compiler infrastructure for multicores

Cetus does automatic parallelization:

- by using data dependence analysis
- by array and scalar privatization
- by reduction variable recognition
- by induction variable substitution

```
#include<stdio.h>
```

```
int main(void)
```

```
{
```

```
    printf("\n Hello World\n");
```

```
    return 0;
```

```
}
```

```
$ gcc main.c
```

```
$ gcc main.c -o main
```

```
# The output of preprocessing stage
```

```
$ gcc -E main.c > main.i
```

```
# Assembly level output
```

```
$gcc -S main.c > main.s
```

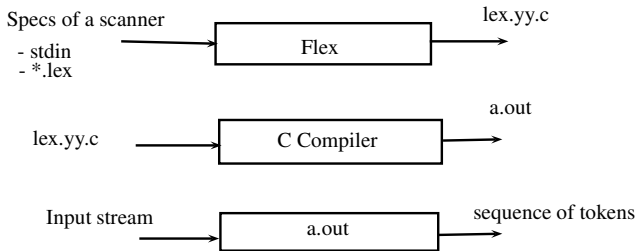
```
#compiled code (without any linking)
$gcc -C main.c
#all the intermediate files using -save-temps function
$ gcc -save-temps main.c
$ ls
a.out main.c main.i main.o main.s
# Print all the executed commands using -V option
$ gcc -Wall -v main.c -o main
Using built-in specs.
COLLECT_GCC=gcc
COLLECT_LTO_WRAPPER=/usr/lib/gcc/i686-linux-gnu/4.6/lto-wr
Target: i686-linux-gnu
Configured with: ../src/con ....
Thread model: posix
gcc version 4.6.3 (Ubuntu/Linaro 4.6.3-1ubuntu5)
```

# lexical analyser generator

```
$ gedit test.lex
/* just like Unix wc */
%{
int chars = 0;
int words = 0;
int lines = 0;
%}
%%
[a-zA-Z]+ { words++; chars += strlen(yytext); }
\n      { chars++; lines++; }
.       { chars++; }
%%
int main(){
yylex();
printf("%d %d %d\n", lines, words, chars);
return 0;
}
```

```
$ flex test.lex  
$ ls  
$ gcc lex.yy.c -lfl ; link to flex library  
$ ./a.out
```

# lexical analyser generator...



- FLEX (Fast LEXical analyzer generator) is a tool for generating scanners.
- First, FLEX reads a specification of a scanner either from an input file \*.lex, or from standard input, and it generates as output a C source file lex.yy.c.
- Then, lex.yy.c is compiled and linked with the “-lfl” library

## lexical analyser generator...

- \*.lex is in the form of pairs of regular expressions and C code.
- lex.yy.c defines a routine yylex() that uses the specification to recognize tokens; a.out is actually the scanner!

Command Sequence:

```
flex sample*.lex
gcc lex.yy.c -lfl
./a.out
```

Input file Format:

```
definitions
%%
rules
%%
user code
```

The definitions section: "name definition"

The rules section: "pattern action"

The user code section: "yylex() routine"

links:

- <http://epaperpress.com/lexandyacc/>
- <http://dinosaur.compilertools.net/>
- <http://dinosaur.compilertools.net/lex/>
- [http://en.wikipedia.org/wiki/History\\_of\\_compiler\\_construction](http://en.wikipedia.org/wiki/History_of_compiler_construction)
- <http://www.drdobbs.com/database/lex-and-yacc/184409830>