

# **Conventional Encryption Message Confidentiality**

# What Is Cryptography?

- Cryptography -- from the Greek for “secret writing” -- is the mathematical “scrambling” of data so that only someone with the necessary *key* can “unscramble” it.
- Cryptography allows secure transmission of private information over insecure channels (for example packet-switched networks).
- Cryptography also allows secure storage of sensitive data on any computer.

# Classical Cryptography: Secret-Key or Symmetric Cryptography

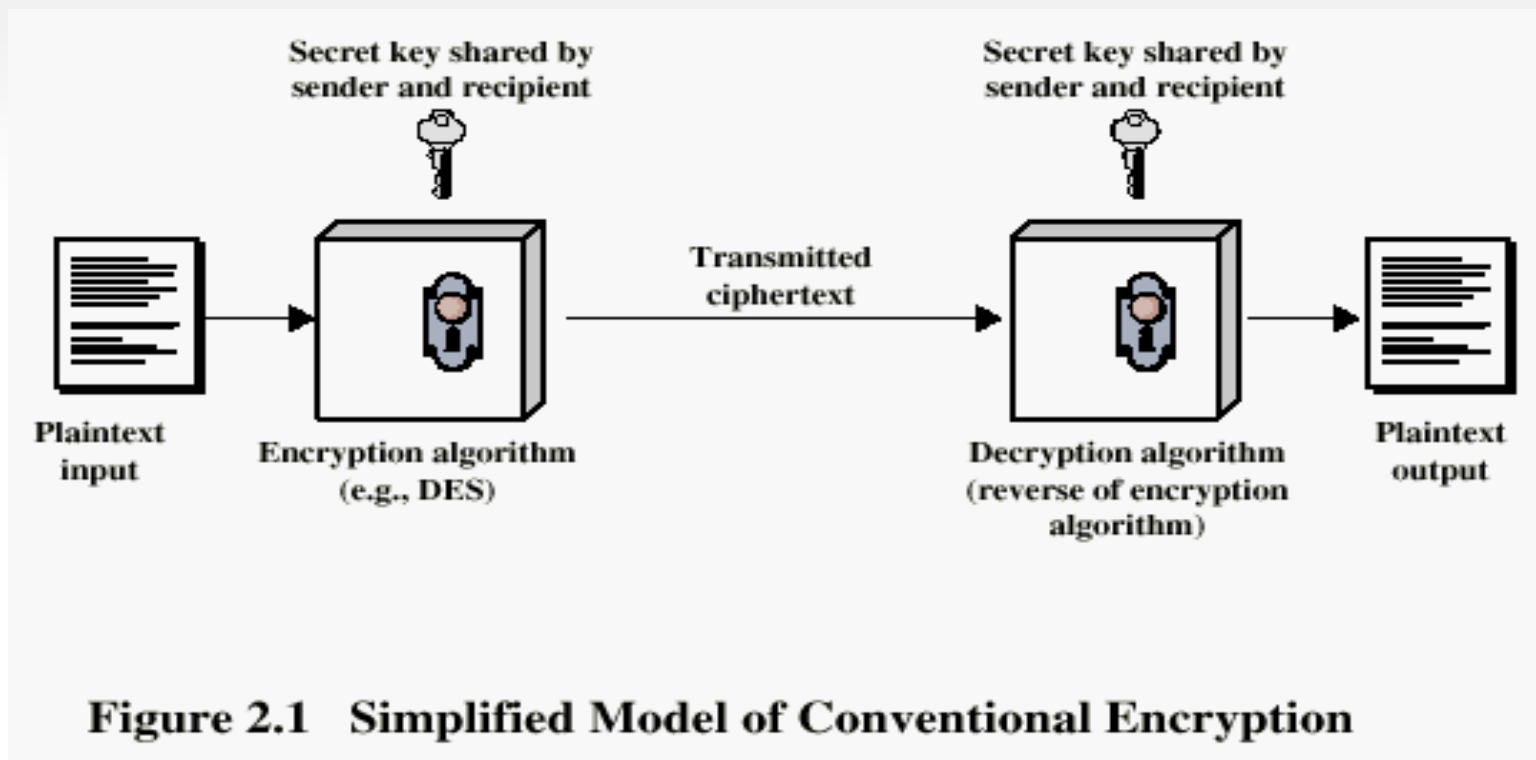
- Alice and Bob agree on an encryption method and a shared *key*.
- Alice uses the key and the encryption method to *encrypt* (or *encipher*) a message and sends it to Bob.
- Bob uses the same key and the related decryption method to *decrypt* (or *decipher*) the message.

# Conventional Encryption Principles

- An encryption scheme has five ingredients:
  - Plaintext
  - Encryption algorithm
  - Secret Key
  - Ciphertext
  - Decryption algorithm
- Security depends on the secrecy of the key, not the secrecy of the algorithm !

# Conventional Encryption Principles

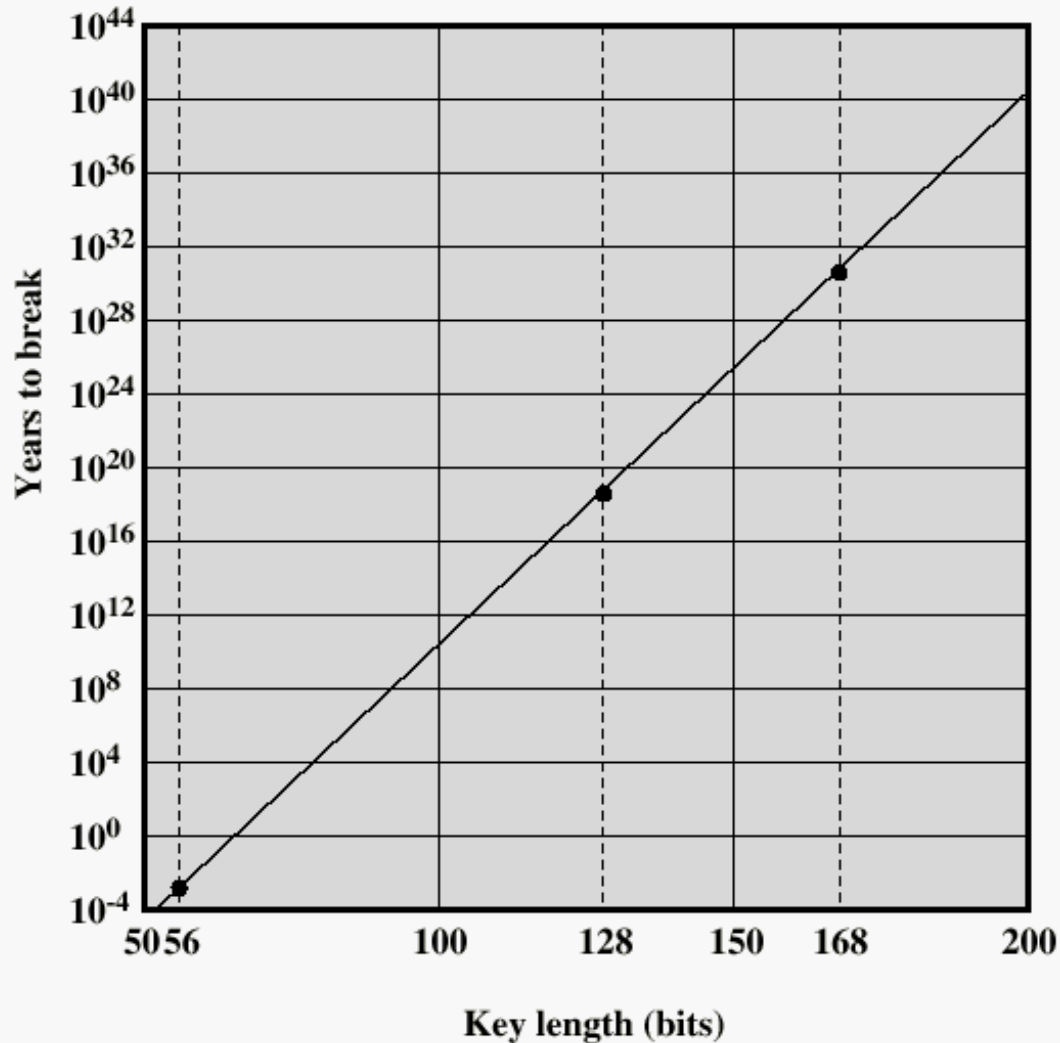
$x$ =plain text,  $k$ = key, ciphertext  $y=E[k,x]$ , decipher:  $x=D[k,y]$



# Average time required for exhaustive key search

Key Size (bits)	Number of Alternative Keys	Time required at $10^6$ Decryption/ $\mu$ s
32	$2^{32} = 4.3 \times 10^9$	2.15 milliseconds
56	$2^{56} = 7.2 \times 10^{16}$	10 hours
128	$2^{128} = 3.4 \times 10^{38}$	$5.4 \times 10^{18}$ years
168	$2^{168} = 3.7 \times 10^{50}$	$5.9 \times 10^{30}$ years

# Time to break a code ( $10^6$ decryptions/ $\mu$ s)



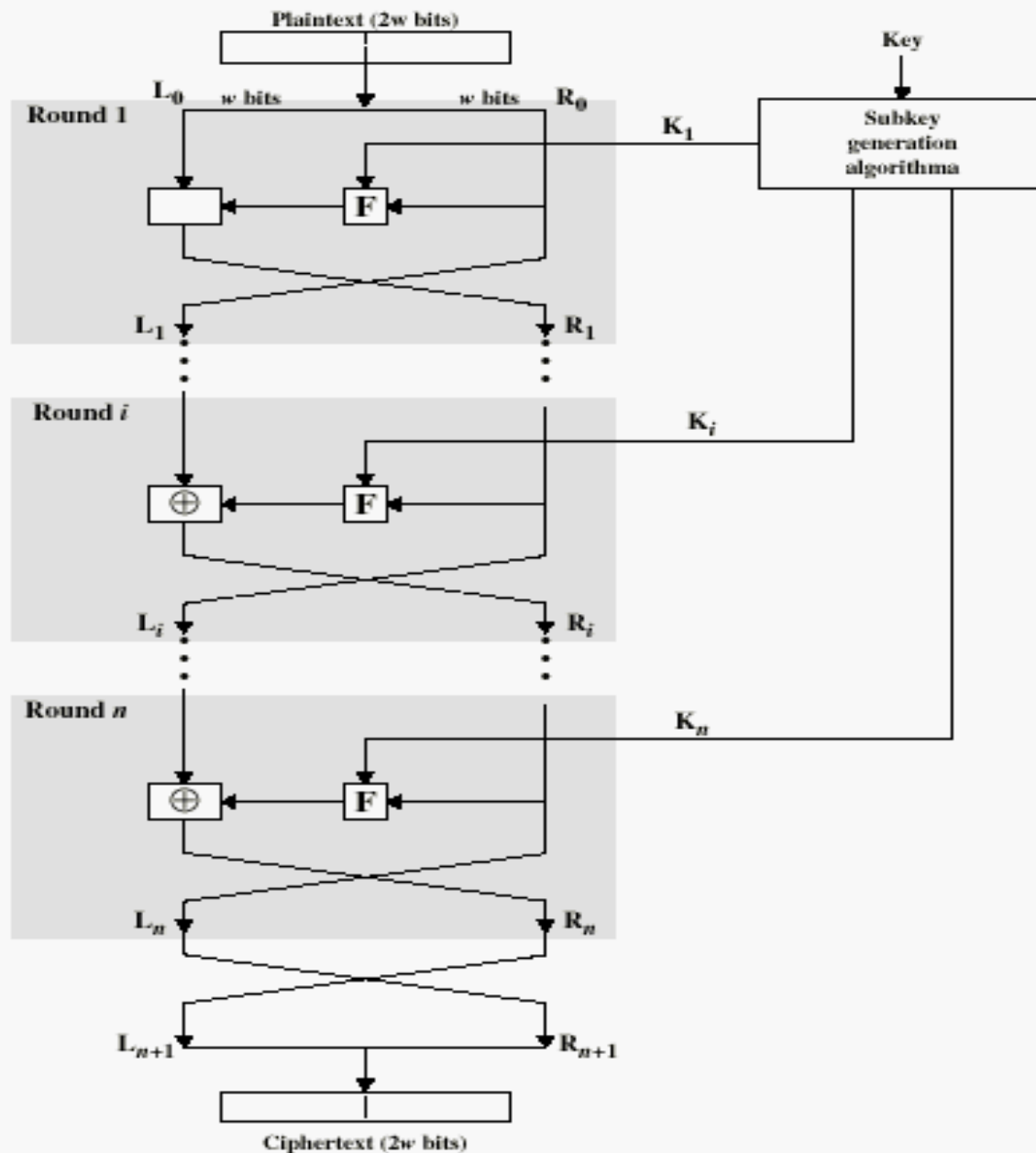
# Feistel Cipher Structure

- Virtually all conventional block encryption algorithms, including DES have a structure first described by Horst Feistel of IBM in 1973
- The realization of a Feistel Network depends on the choice of the following parameters and design features (see next slide):



# Feistel Cipher Structure

- **Block size:** larger block sizes mean greater security
- **Key Size:** larger key size means greater security
- **Number of rounds:** multiple rounds offer increasing security
- **Subkey generation algorithm:** greater complexity will lead to greater difficulty of cryptanalysis.
- **Fast software encryption/decryption:** the speed of execution of the algorithm becomes a concern



**Figure 2.2 Classical Feistel Network**

# DES

- The overall processing at each iteration:
  - $L_i = R_{i-1}$
  - $R_i = L_{i-1} \oplus F(R_{i-1}, K_i)$
- Concerns about:
  - The algorithm and the key length (56-bits)

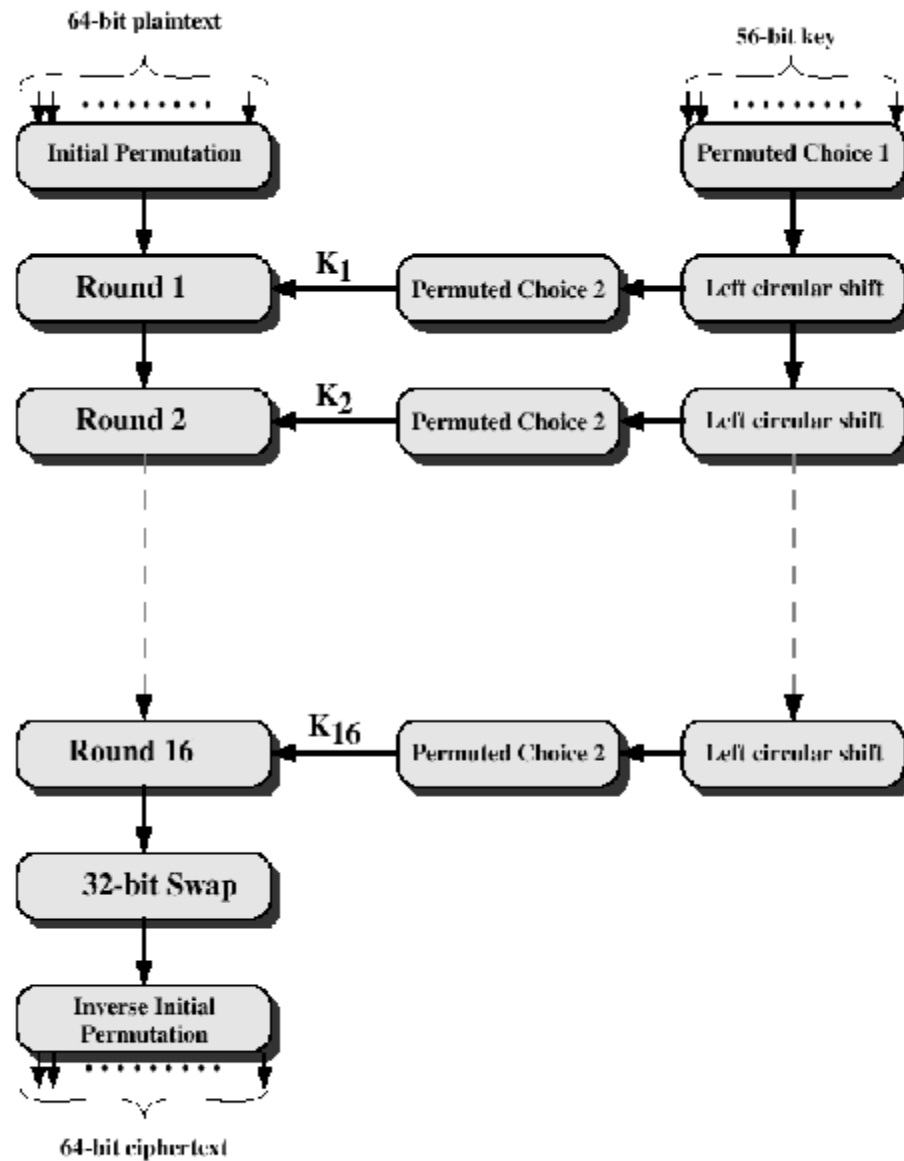


Figure 2.3 General Depiction of DES Encryption Algorithm

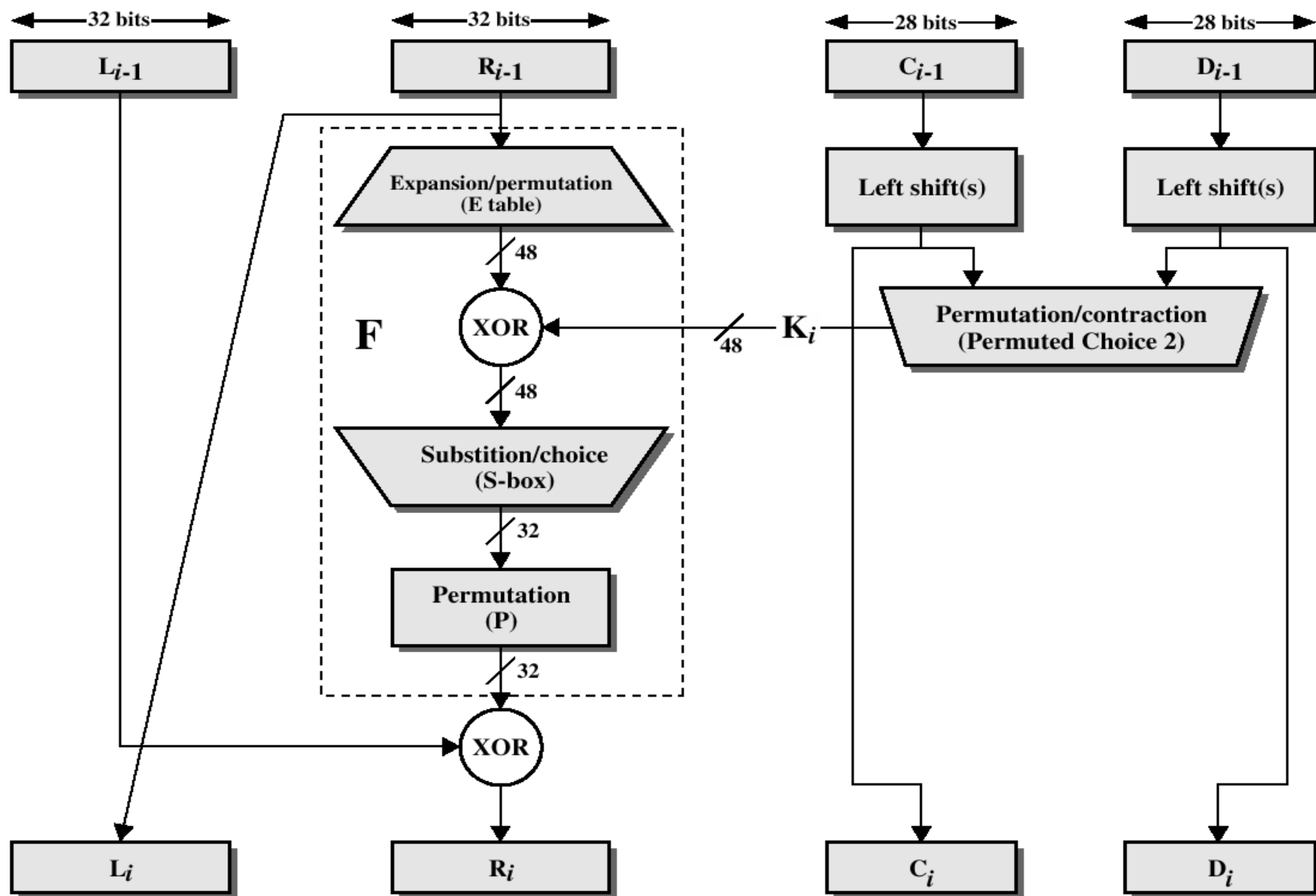
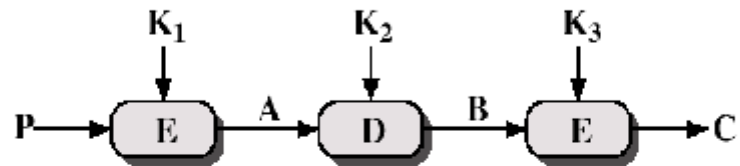
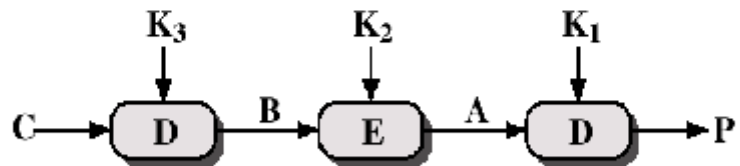


Figure 2.4 Single Round of DES Algorithm

# Triple DEA



(a) Encryption



(b) Decryption

Figure 2.6 Triple DEA

# Cipher Block Modes of Operation

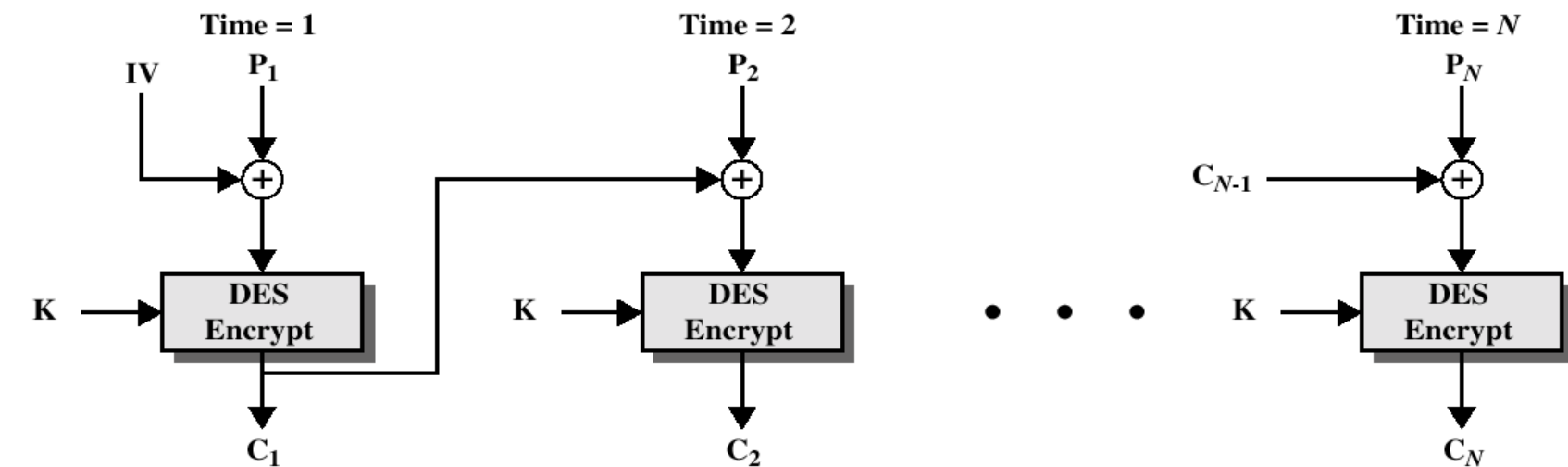
- Cipher Block Chaining Mode (CBC)
  - The input to the encryption algorithm is the XOR of the current plaintext block and the preceding ciphertext block.
  - Repeating pattern of 64-bits are not exposed

$$C_i = E_k[C_{i-1} \oplus P_i]$$

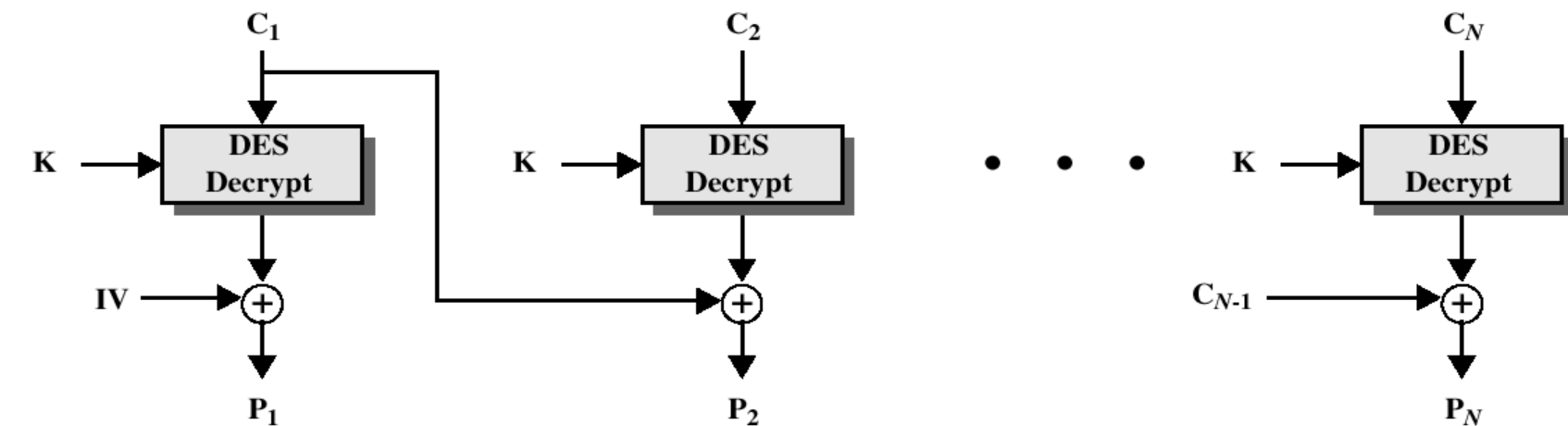
$$D_K[C_i] = D_K[E_K(C_{i-1} \oplus P_i)]$$

$$D_K[C_i] = (C_{i-1} \oplus P_i)$$

$$C_{i-1} \oplus D_K[C_i] = C_{i-1} \oplus C_{i-1} \oplus P_i = P_i$$



(a) Encryption



(b) Decryption

Figure 2.7 Cipher Block Chaining (CBC) Mode



# Advantages of Classical Cryptography

- There are some very fast classical encryption (and decryption) algorithms
- Since the speed of a method varies with the length of the key, faster algorithms allow one to use longer key values.
- Larger key values make it harder to guess the key value -- and break the code -- by brute force.

# Disadvantages of Classical Cryptography

- *Requires secure transmission of key value*
- Requires a separate key for each group of people that wishes to exchange encrypted messages (readable by any group member)
- For example, to have a separate key for each pair of people, 100 people would need 4950 different keys  $n(n-1)$ .

# Clasification of Cryptography

- The number of keys used
  - symmetric (single key)
  - asymmetric (two-keys, or public-key encryption)

# Conventional Encryption Algorithms

- Data Encryption Standard (DES)
  - The most widely used encryption scheme
  - The algorithm is referred to the Data Encryption Algorithm (DEA)
  - DES is a block cipher
  - The plaintext is processed in 64-bit blocks
  - The key is 56-bits in length

# Triple DEA

- Use three keys and three executions of the DES algorithm (encrypt-decrypt-encrypt)

$$C = E_{K3}[D_{K2}[E_{K1}[P]]]$$

- $C$  = ciphertext
  - $P$  = Plaintext
  - $EK[X]$  = encryption of  $X$  using key  $K$
  - $DK[Y]$  = decryption of  $Y$  using key  $K$
- Effective key length of 168 bits

# Location of Encryption Device

- **Link encryption:**
  - A lot of encryption devices
  - High level of security
  - Decrypt each packet at every switch
- **End-to-end encryption**
  - The source encrypt and the receiver decrypts
  - Payload encrypted
  - Header in the clear
- **High Security:** Both link and end-to-end encryption are needed

## Secret key Distribution:

- The two parties must share they secret key
  1. generated by one party and manually deleivered
  2. generated by one and delivered using shared secret
  3. created by thid party and deliver through secured channel
- 1-2: for small organizations
- number of pairs for n users:  $n*(n-1)/2$
- KDC can play the role of key distribution

## Algorithms:

1.  $U_A$  issues req. to kdc for session key for session with user  $U_B$

req. = identity of  $U_A$ , identity of connection  $ID_A$  (=time stamp, counter or rand. no., protocol ref.)

2. kdc creates session key( $K_s$ ), responds with msg encrypted using a pre-established secret key  $K_A$  between  $U_A$  and kdc.

msg = session key  $K_s$  + original req copy + sub msg.  
encrypted using preshared key  $K_B$  between kdc and  $U_B$  (sub msg= $K_s$ , id. of  $U_A$ )



Now  $U_A$  can verify that original req. was not altered.

3.  $U_A$  sends msg to  $U_B$ .

msg = its identity, session id  $ID_A$ , encrypted sub message of kdc to be relayed to  $U_B$ , encryption of  $ID_A$

4. On arrival  $U_B$  can check identity of  $U_A$ , and integrity of Ks.

(At this point one can deduce that secret key has been securely delivered to  $U_A$  and  $U_B$ .)

Now the message transfer can go on between  $U_A$  and  $U_B$ .